This homework will require you to translate between English descriptions of typing rules and formal inference rules. For example, we have the following typing rule for product types:

If x has type σ and y has type τ , then (x,y) has type $\sigma * \tau$

which is represented by the formal inference rule

$$\frac{\Gamma \vdash \mathbf{x} : \sigma \quad \Gamma \vdash \mathbf{y} : \tau}{\Gamma \vdash (\mathbf{x}, \mathbf{y}) : \sigma * \tau}$$

So, we'll start by giving you either the English description or the formal rule, and asking you to produce the other. Then, we'll move on to giving multiple rules governing a more complex type.

If you're trying to LATEX this, an easy way to typeset this is by using the proof package (i.e. usepackage{proof}). Then, you can produce the above inference rule with

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{\Gamma\vdash\texttt{(x,y)}:\sigma * \tau}
{\Gamma\vdash\texttt{x}:\sigma & \Gamma\vdash\texttt{y}:\tau}
```

Question 1 (Required)

What does the following rule say in English?

$$\frac{\Gamma \vdash \mathtt{f}: \sigma \to \tau \quad \Gamma \vdash \mathtt{t}: \sigma}{\Gamma \vdash (\mathtt{f} \ \mathtt{t}): \tau}$$

Question 2 (Required)

Write the following rule as a formal inference rule.

If, in context Γ , **b** is of type **bool**, and both **e1** and **e2** are of type τ (in Γ), then the expression **if b then e1 else e2** is of type τ

Question 3 (Required)

For any type τ , we define the type τ option by the following rules:

- In any context, NONE : τ option
- If x is of type τ in context Γ , SOME(x) is of type τ option in context Γ
- If
 - e1 : σ in context Γ
 - $e2: \sigma \text{ in context } \Gamma, \mathbf{x}: \tau$
 - opt : τ option in context Γ

then

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(case opt of NONE => e1 | (SOME x) => e2)
```

has type σ in context Γ

Write these rules as formal inference rules

Question 4 (Optional)

Write inference rules for the type **bool** of booleans. You should include the values of type bool, rule(s) for getting booleans from other types (e.g. if x and y are of the same type, then x=y is a boolean¹), and rules for using booleans (e.g. what's an appropriate rule for if...then...else expressions?).

 $^{^1\}mathrm{Don't}$ worry about equality types

Question 5 (Optional)

Come up with an inference rule governing how to typecheck let...in...end expressions with only a single declaration between the let and the in, e.g.

let
 val x = e
in
 y
end